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Software Engineering Department

Braude College of Engineering

Capstone Project Phase B

**Path Planning Via BMC of Hyper-Properties**

**23-2-R-5**

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Link to the project’s GitHub repository:

[GitHub - BMC Path Planning](https://github.com/danielbal21/BMC_Path_Planning)

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**Abstract**

*Model Checking* is a verification technique that explores a state space of a system to determine if it obeys a specification of its intended behavior. *Bounded Model Checking (BMC)* iteratively explores the state space of a system, exploring all possible states up to a certain bound[2], in search for a counterexample, which is a system run that violates a required specification. *hyper-properties* are used to describe the behavior of a combination of systems. This research project presents an application of BMC for hyper-properties to find a path for an agent that does not intersect with dynamic counteragents on a grid. This is achieved by translating the systems (agent and counter-agents systems) into Boolean Formulas and using *SAT-Solvers* for finding a path if exists.

we present the algorithm, data structures, software design and implementation while providing elaborate explanations of the processes, tools and evaluation plans.

**1. Introduction**

Pathfinding on a grid with moving agents is a problem of determining whether there exists a valid path between two points on a grid-based map.[6]

The pathfinding process is done on a 2D grid-based environment, where the grid size can vary. Multiple agents are moving on the grid in an indeterministic fashion.

There is one agent (main agent) which always starts in point and its goal is to reach point . in our case is the bottom left and is the top right of the grid.

This problem is often encountered in scenarios where multiple entities, such as robots, vehicles, or characters in a video game, need to navigate a shared environment without colliding with each other.

The pathfinding algorithm needs to consider the dynamic nature of the agents' movements, anticipating their positions and avoiding them accordingly. It involves analyzing the current and predicted positions of the agents and calculating a path that avoids potential collisions or intersections. Common algorithms may need to continuously update the path as the positions of the agents change over time.

The objective is to determine if a feasible path exists, rather than finding the shortest or most optimal path. The focus is on avoiding the agents and ensuring safe navigation rather than minimizing distance or cost.

In our project the environment is a von Neumann grid; A von Neumann grid is a 2D grid structure where each cell is connected to its four neighboring cells: up, down, left, and right. The von Neumann grid offers simpler connectivity compared to the more intricate Moore grid, as it allows movement or connection in only four cardinal directions.  
however, a fifth direction which is to stay in place is also possible for the agent.

For BMC to be applied into pathfinding problems, the problem needs to be modeled in a way that can eventually be represented as a *satisfiability problem*.[1]

Given a grid the size of ( a main agent and a counter-agent, we describe the main agent and the counter-agents behavior by using *Kripke structures* and respectively.  
Each path of can be represented by a sequence of states and transitions from the initial state to a final state (if exists).

Each state of represents where how are the counter-agents positioned across the grid.  
Each state of and is an image of the current position of all agents in the system on the grid.

The transitions describe possible future images that can exist at the next iteration.

The green robot can move up or right, one step in every iteration while the red robot moves vertically on an infinite course, one step at the time, The green robot seeks to reach the top right corner (figure 1).

Graphical user interface

Description automatically generated with medium confidence

Figure 1 – The Idea

The ‘green robot’ is referred to as the *main agent* while red robots are referred to as *counteragents.*

a simple way to answer the question “is there a path of that does not intersect with is to explore all possible paths of and all possible paths of . However, this solution is highly inefficient since it involves scanning all possible paths of both systems.

This can be modeled by using Kripke Structures (figure 2).

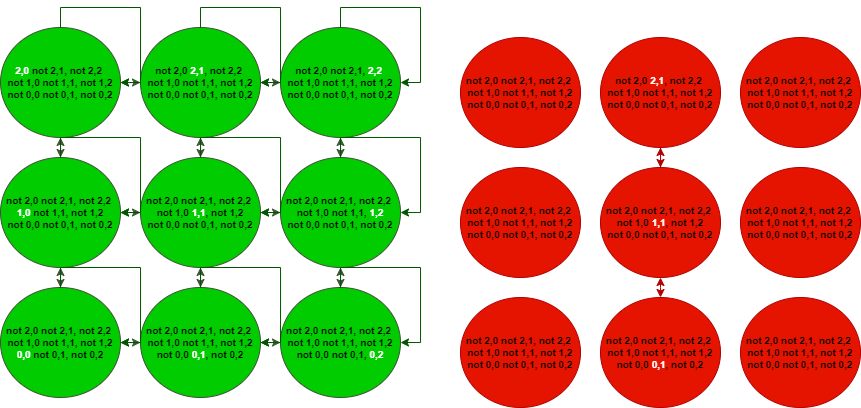
**

Figure 2 – Complete Kripke Structures

Where for each state the agent is either in the coordinates or not (green system) and for each counteragent the is not negated if the counter agent is in the position (red system). these forms of structures can be converted into Boolean formulas which then can be used with BMC to find a safe path.

In this project we aim to check this approach efficiency and implement the system’s creation process, translate the systems into Boolean formulas, feed them to a SAT-Solver and extract the path from the solution (if exists). Finally, we ran experiments on various systems that differ in size and agent count to evaluate the algorithm efficiency and limitations.

# 2. Background and Related Work

**2.1 Kripke Structures**

A Kripke Structure is a variation of a transition system, used in model checking to represent the behavior of a system.  
It consists of a graph where each node represents a state of the system, and the edges represent state transitions.  
Each node in the structure holds a set of properties.[4]  
**Definition 1.**  Let be a set of atomic propositions representing the properties. A Kripke Structure over is a tuple , where is a set of states, is a set of initial states, is the transition relation of the structure and is the labeling function.  
Definition 2. A path in a Kripke Structure is an infinite sequence of states such that and for all .  
 for each node in the sequence, there must be a transition from to .

**Definition 3.** The *Language* of a Kripke Structure is the set of all possible paths.

For example, let us consider the system:  
   
a path of the system is ( and are a tuple in the relation)

**Motivation**To further clarify why we should use a Kripke structure we will consider the following analogy:   
Let's define a counter agent system placed on an grid. Each step in the system is triggered by pressing a "step button." Upon each button press, all the counter agents in the system make a move.

To capture the system's behavior, we use a camera to take a picture after each button press. However, due to the nondeterminism of the agents' movements, each agent can potentially move in multiple directions. This results in multiple possible pictures for each button press.

We associate each picture with the picture it originated from, forming a connection between states in the system. This collection of connected states and their associated pictures forms a Kripke structure that accurately describes the system’s behavior.  
By exploring the "gallery" of pictures connected through their originators, we can identify and create a safe path within the counter agent system. This path represents a sequence of states that guarantees the absence of collisions or unsafe situations among the agents.[7]

**2.2 SAT Problem**

**Boolean Formula**a Boolean formula is a finite string which is constructed from:

* Variables:
* Boolean operators:
* Parenthesis:

Each of the variables in a formula can have the value **true** or **false**  
a formula is the result **true** or **false**  under the *assignment*

For example, let us consider the formulaand the assignments  
 ,   
the value of under can be written as and is evaluated to  
 but

Boolean formulas are used to show Satisfiability, Unsatisfiability, Validity, Contradiction and Tautology:  
Let be a Boolean formula.

- is *valid* for every assignment is evaluated to **true** or **false.**

example:

- is *satisfiable* there exists an assignment which is evaluated to **true.**

example:

- is *unsatisfiable (contradiction)* for every assignment is evaluated to **false.**

example:

- is a tautology for every assignment is evaluated to **true**. example:

The Boolean satisfiability problem (SAT) is a problem of finding an assignment to a set of Boolean variables such that a Boolean formula evaluates to ‘True’ under . Given that SAT is the first problem that was proved to be [NP-complete](https://en.wikipedia.org/wiki/NP-complete), completing such a task is not straightforward, since there isn’t any known polynomial algorithm, that for a given SAT formula can determine whether its satisfiable or not. [5]

## **2.3 SAT Solver**

*SAT Solvers* are programs that aim to solve the SAT problem. For a given Boolean formula , a SAT Solver outputs whether is satisfiable, and if so, returns the set of Boolean variables that satisfies . Since SAT is NP-C the solvers improve runtime by using sophisticated heuristics that enable them to solve very large Boolean formulas. Most SAT solvers include time outs, so they will stop attempting to solve the problem if they haven’t found a solution in a reasonable time and will output ‘unknown’[5]. For example, consider the following two Boolean formulas: ,   
Giving that is satisfiable, feeding it to a SAT solver would result in an output such as:   
But, feeding to a SAT solver would result in the output: .

In our project we use *Z3* as our solver.

Solvers like Z3 can handle Boolean formulas of varying complexity and size, from relatively small and simple formulas to larger and more intricate ones. While there is not a strict limit on the scale of formulas, solver performance may degrade as the complexity and size of the formulas increase, potentially leading to longer solving times or resource constraints.

## **3. Expected Achievements**

In our project, we expect to achieve a working efficient path-finding algorithm on an grid involving a primary agent going from start to end while avoiding moving counter-agents. The algorithm is based on BMC using encoded Kripke Structures generated from the agent and counteragents systems.

We also expect to generate an operational GUI with a visualization layer  
and various general abilities (System generation, import, export, benchmarking, and analyzers).  
another part of our project involves benchmarking the algorithm for various input systems of different sizes and complexity, and to test the algorithm’s ability to provide a solution in reasonable time.

# 4. The Solution

# 4.1 Learning Stage (Processing Pipeline)

The processing pipeline consists of five stages.

1. **Input** – At the outset, the pipeline accepts input in the form of system specifications, including grid dimensions, counteragent configurations, and any relevant parameters necessary for subsequent stages.
2. **Grid Cells and Robots** – In this phase, the system's architecture is abstracted into a structured representation, delineating grid cells and identifying the presence or absence of counteragents within each cell. This step establishes the foundational framework upon which subsequent analysis built.
3. **Flattening** – During flattening, the evolving dynamics of the system are distilled into a temporal sequence of snapshots, each capturing the state of grid occupancy at a specific point in time. This process facilitates the temporal abstraction of the system, essential for modeling its temporal behavior.
4. **Kripke Structure** – Transitioning to a more formal representation, the pipeline constructs a Kripke structure to encapsulate the system's evolution over time. Here, nodes correspond to discrete time points, while transitions between nodes capture the temporal progression of the system. This structural transformation enables rigorous analysis of temporal properties and behavior.
5. **Boolean Formula** – Finally, the system's dynamics, constraints, and objectives are encoded into a logical Boolean formula. By expressing system properties in a formal logic framework, the pipeline facilitates the application of automated solvers to ascertain solutions that meet specified criteria. This logical formulation provides a powerful means to explore, analyze, and reason about the system's behavior and potential outcomes.



Figure 3 – Processing Pipeline

# 4.2 System Generator

The automatic generation of a system involves creating a grid environment populated with counteragents that move according to certain parameters.

*Grid Size*  - This parameter defines the size of the grid environment, specifically an grid. It determines the dimensions of the space in which the counteragents will navigate.

*Number of Counter Agents* - This parameter specifies the number of counteragents to be generated in the grid environment. It controls the density of counteragents within the space and affects the complexity of the pathfinding problem.

*Stay Chance\** - This parameter represents the probability that a counteragent will choose to stay in place rather than move to an adjacent grid cell during each time step. It influences the likelihood of counteragents remaining stationary, adding variability to their movement behavior.

*Stray Radius* - This parameter defines the maximum distance that a counteragent can move away from its initial position in any direction. It constrains the extent to which counteragents can wander from their starting locations, affecting the overall spread of counteragent positions within the grid.

\* The *Stay Chance* parameter is critical as it dictates whether a counteragent remains stationary in a grid cell or moves. This decision has far-reaching consequences as staying in a cell permanently blocks it and all potential paths (potentially) originating from it, significantly impacting the grid's accessibility and path-solving feasibility.

The counteragent system generator utilizes an algorithmic framework, integrating probabilistic decision-making and recursive traversal methodologies to emulate the behavior of counteragents within a grid-based environment. Initially, the algorithm randomizes the placement of the counteragent within the grid, introducing stochasticity to its initial position and enhancing the diversity of generated scenarios.

Subsequently, at each stage of traversal, the counteragent evaluates potential movements based on predefined probabilities. Notably, these probabilities, except for the decision to remain stationary, are uniformly distributed at 50%, ensuring equitable exploration of directional options. This balanced approach contributes to a more nuanced representation of counteragent behavior, facilitating the generation of diverse movement patterns.

Upon selecting a movement direction, the counteragent calculates its trajectory from the new position, accounting for factors such as the distance traveled from the origin. Should the calculated travel distance exceed a specified stray radius, the counteragent is compelled to cease movement and maintain its current location.

Furthermore, the algorithm dynamically adapts to evolving scenarios, The detection of revisited grid cells serves as a crucial indicator of path completion, enabling the identification of closed loops and the evaluation of path efficiency.

The outcome of this generation process yields a refined input format, aligning with the subsequent stage of *Grid Cells and Robots* in the processing pipeline.

# 4.3 Flattening the Counteragents

The Flattening algorithm is an essential intermediary step in transforming the grid-based representation of a counteragent system into a Kripke structure. This algorithm operates by constructing a timeline where each step corresponds to a matrix representing all potential occupancies of counteragents at that moment in time.

The Flattening algorithm serves a critical purpose in streamlining the representation of a counteragent system into a more manageable form. Without this algorithm, constructing a Kripke structure that considers each counter-agent path individually would lead to an explosion in the number of atomic propositions, significantly lengthening the boolean formula. Additionally, the complexity arising from analyzing each counteragent and creating infinitely many realities where each counteragent can move freely introduces significant computational burdens. By flattening the grid from multi-counteragent configurations to single-occupancy matrices, the algorithm sacrifices individuality for simplicity, resulting in a more tractable model.[3]

At the onset of the algorithm, an initial node is established to kickstart the timeline. This initial node encapsulates the starting positions of all counteragents within the system. Each counteragent’s position is recorded, and a record of the counteragent positions is maintained for subsequent steps.

The algorithm iterates through successive steps in the timeline until a closed loop is detected, signifying the completion of the process. During each iteration, the potential whereabouts of each counteragent in the next step are calculated based on their current positions and movement capabilities. These potential positions are examined to determine the possibility of movement in each direction.

The algorithm's convergence implies that it will eventually reach a state where it finds a movement pattern for all counteragents that has been encountered previously, thereby concluding the process. This convergence ensures that the algorithm exhaustively explores all possible movement scenarios for each counteragent.

During the process, new nodes are added to the structure to represent each successive step in the timeline. Relations between nodes are established to denote the transition from one step to the next.

The algorithm's convergence is essential due to a key constraint: each counteragent must demonstrate cyclic behavior, represented by closed loops in their movements. As a result, the algorithm strives to identify a cyclic timeline encompassing all potential cycles for each counteragent, marking the completion of the process.

In figure 4, counter agent 1 and counter agent 2 with a motion cycle of 1 and 2 or 3 respectively are flattened into a potential occupancy timeline with motion cycle of 3.

Where a Kripke structure is generated to represent the occupancy indicators at each timestep.

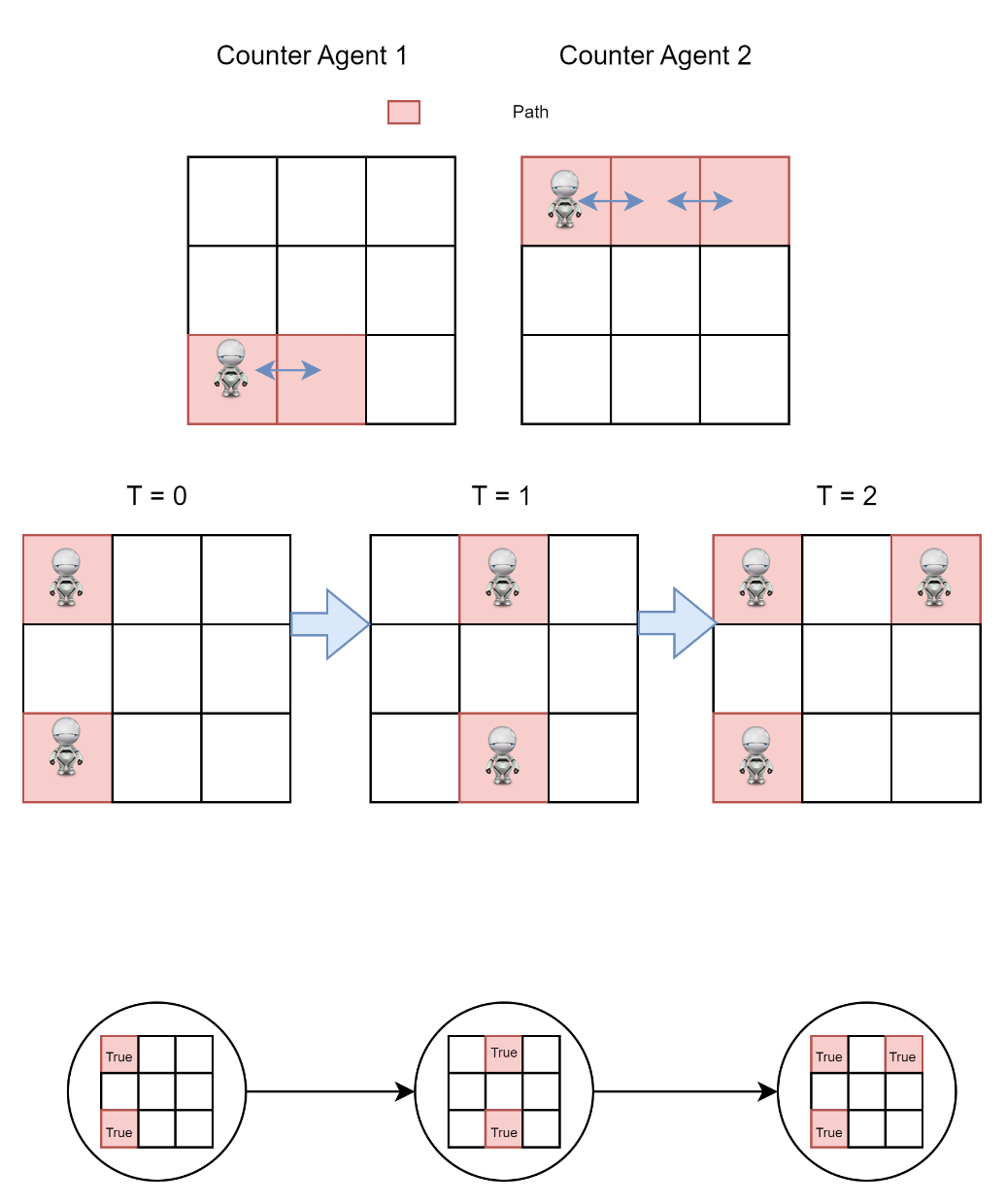


Figure 4 – Flattening Process

# 4.4 Algorithm

The algorithm aims to find a valid path for an agent moving on a grid represented by a Kripke structure, ensuring that it avoids collision with obstacles represented by another system (M2).

**This algorithm takes place after the counteragents Kripke structure is formed.**

The grid is represented as a set of nodes in a Kripke structure (M1), where each node represents a state of the grid.

Each state contains a Boolean matrix representing the *occupancy* of each cell in the grid. A *True* value indicates the presence of an obstacle or the counteragent itself in that cell, while *False* indicates an empty cell.

The movement of all agents is represented by transitioning between states in the Kripke structure. Each transition corresponds to moving the agent to an adjacent cell in the grid, including staying in the same cell.

**Constraints for Valid Movement:**The algorithm ensures that the main agent moves in a valid manner by imposing constraints on its movement,allowing the main agent to move only to adjacent cells. This ensures that the path accepts grid movement laws.Additional constraints are applied to ensure that the main agent stays within the bounds of the grid.

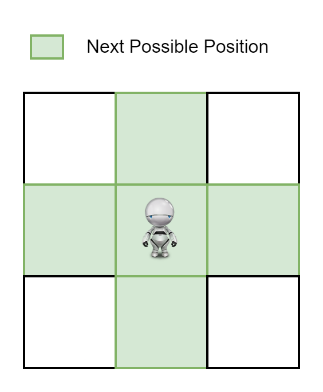


Figure 5 – Valid Movement Constraint

**Constraints for Path Finding:**

Constraints are formulated to ensure that there exists a path from the initial state to the final state while satisfying the movement constraints defined earlier.

The path is constructed using Boolean variables to represent the movement of the main agent at each time step.

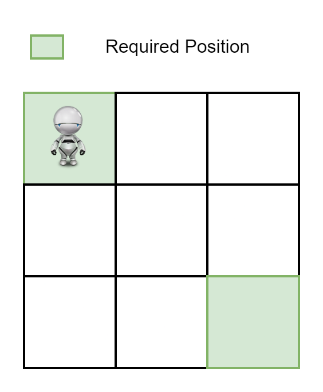


Figure 6 – Path Finding Constraint

**Safety Constraints:**To ensure safety, constraints are defined to prevent the main agent from colliding with obstacles represented by the second system (M2).

These constraints ensure that at each time step, the main agent does not occupy cells that are occupied by obstacles.

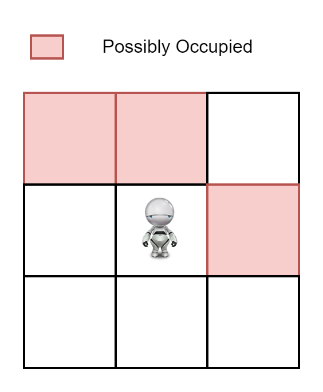


Figure 7 – Safety Constraints

**Solving the Constraints:**The formulated constraints are provided to a solver (such as Z3) to find a solution that satisfies all constraints.

If a solution is found, it represents a valid path for the main agent to traverse from the initial state to the final state without colliding with obstacles.  
If no solution is found, it indicates that no valid path exists, considering the defined constraints.

# 4.5 Formula Creation Process

The solution involves formulating the problem as a set of logical constraints using Boolean variables. This process ensures that we capture the essential requirements for finding a valid path for the main agent while avoiding collisions with counter agents and staying within the grid boundaries. Each constraint corresponds to a specific condition that the solution must satisfy, and these constraints are combined using logical operators to form a comprehensive problem formulation.

The base boolean variables, denoted as represents the presence or absence of the main agent in each cell of the grid at each time step and represents the row and column respectively. These variables are the building blocks of our problem formulation and are used to encode the state of the system at each point in time.

The first formula *Initial Position* ensures that the main agent starts at the initial position which is fixed to .

The second formula *Final Position* ensures that the main agent ends at the final position which is fixed to where is the grid size and is path length.

The third formula *Valid Steps* ensures that the main agent can only move to neighboring cells or stay in its current cell at each time step, thus defining valid movements options.

\* Note that there is a grid boundary checking that is not shown in for simplicity concerns.

The fourth formula *Safety* ( ensures that the main agent cannot occupy cells with counter agents at any time step, thereby ensuring safety.

\* Where is the counter agents system occupancy indicator at

The fifth formula *Single Path* ensures that the main agent follows a single path by allowing at most one occupied cell per time step, preventing multiple simultaneous movements.

The overall boolean formula to represent the problem is the conjunction (*AND*) of all the individual constraints:

This formula represents the complete set of constraints necessary to ensure the existence of a valid path for the main agent from the initial state to the final state, while avoiding collisions with counter agents and adhering to the rules of movement within the grid.

# 4.6 Bounded Model Checking (BMC)

Formulas are generated for a specified path length, denoted as . If a solver determines a formula to be unsatisfiable, a formula for a longer path is generated and another attempt is made to find a solution. Path length is increased until a predefined bound is reached or a solution is found. If the bound is reached without finding a valid path, the algorithm concludes with the determination that 'no path exists'.

# 4.7 Path Reconstruction Process

Reconstructing the path involves deriving the sequence of grid states traversed by the main agent from the initial to the final state, based on the satisfying assignment acquired from solving the boolean formula.

Let represent the satisfying assignment obtained from the solver, where denotes the Boolean value assigned to the variable at time step and grid position

To reconstruct the path, we iterate through the satisfying assignment and trace the movement of the main agent. Starting from the initial state we follow the sequence of transitions dictated by the satisfying assignment:

Here, denotes the row and column indices of the grid position at time step , as dictated by the satisfying assignment .

By adhering to this process, we systematically reconstruct the trajectory of the main agent, affirming its adherence to the permissible movements delineated in the satisfying assignment while navigating from the initial to the final state of the grid.

# 4.8 Software Development

The software architecture is structured into three distinct layers: Models, Services, and UI. The UI layer exclusively encompasses the user interface elements, facilitating interaction with the end user. This layer communicates with the Services layer, which is responsible for executing backend operations and processing data. The exchange of information between the Services layer and the UI layer occurs through the utilization of data types defined within the Models layer.

The analyzer is an external component, hooking into the backend services and providing automation for analyzing and testing the algorithm.

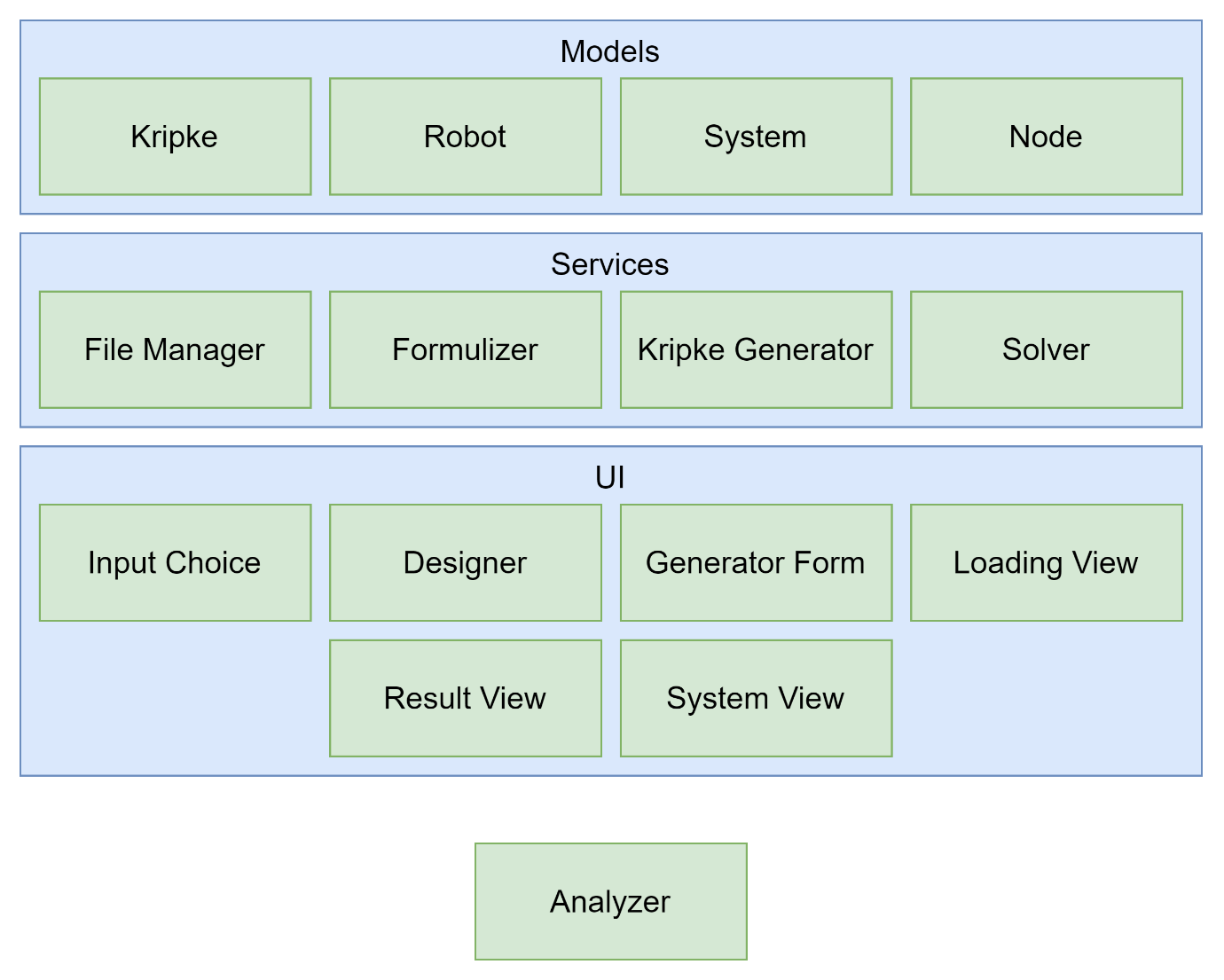


Figure 8 – Top Level Architecture

**User Interface**

input choice refers to the primary menu interface where users select their preferred method of loading the system.

A screenshot of a computer

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Figure 9 – Input Choice Menu

The Designer tool allows users to construct systems interactively, enabling them to define grid dimensions, incorporate counteragents, and designate available movements (up, down, left, right, stay) for each grid cell. Additionally, users can specify the initial cell for each counteragent. As users finalize their grid designs, the Designer tool conducts checks to identify any undefined cells or open loops, ensuring the completeness and coherence of the system.

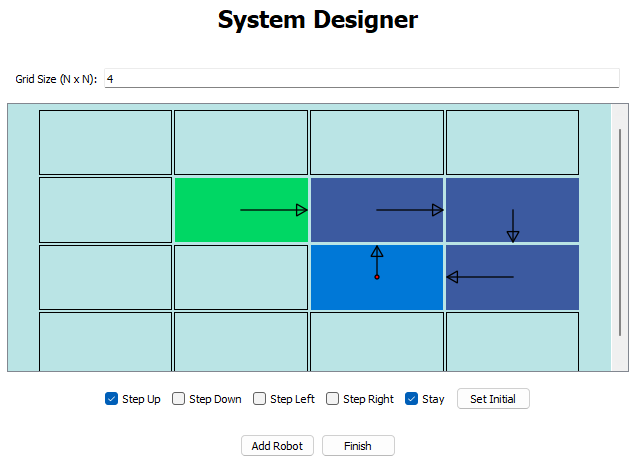


Figure 10 – System Designer

The Generator Form facilitates the automated generation of systems, encompassing all parameters outlined in the system generator section. Users input these parameters into the form, where input validation occurs to ensure the integrity of the data. If invalid data is provided, the form notifies users of any errors, promoting accurate and reliable system generation.

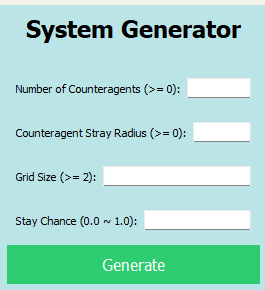


Figure 11 – Generator Form

The System View interface is presented upon successful loading of a system, prompting users to specify runtime parameters. Within this window, users have the option to display a plot showcasing the Kripke structure of the counteragents system. By selecting individual states within the Kripke structure, users can view the corresponding occupancy matrix generated during the Flattening stage. Additionally, the System View interface allows users to save the system into a file, enhancing accessibility and data management capabilities.

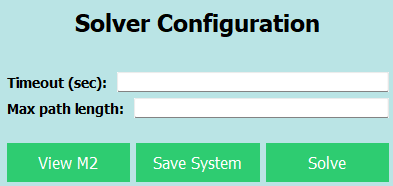


Figure 12 – System View

A screenshot of a computer

Description automatically generated

Figure 13 – Kripke Viewer

The Loading View interface is displayed during the solution phase, providing users with real-time updates on the progress of the solver. It conveys information about the number of iterations attempted out of the maximum path length specified in the solver configuration. Additionally, the Loading View interface indicates the duration elapsed since the solution attempt commenced, enabling users to monitor the solver's performance and estimate remaining processing time.

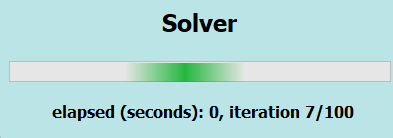


Figure 14 – Loading View

The Result View interface presents the outcome of the solution phase, providing users with essential information regarding the success or failure of the solution process. In the event of a successful solution, the Result View displays the path taken in an animation, illustrating how it navigates to avoid counter agents. Additionally, the interface includes time-point markers for each step in the animation, enabling users to track the progression of the solution. It also calculates and presents the average time per iteration and the total duration of the solution process, offering insights into the efficiency of the solver. In cases where no solution is found, a corresponding message is displayed to inform users of the unsuccessful attempt. Similarly, if the timeout threshold is reached without finding a solution, the interface notifies users, accordingly, ensuring transparency regarding the solver's limitations and constraints.

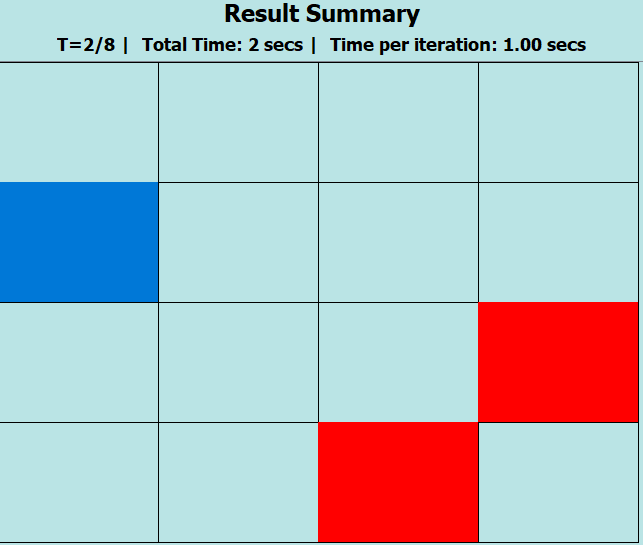


Figure 15 – Result View (Solution Found)

**Services**

*The File Manager* module handles the serialization and deserialization of Kripke objects, facilitating the storage and retrieval of system data. It is responsible for converting Kripke objects into binary data for storage purposes and restoring them back into their original form when needed.  
Loading and Saving the objects is taken/inserted to the Kripke Structure stage in the pipeline.

*The Formulizer* module orchestrates the formulation of logical expressions derived from the Kripke Structure. Its pivotal role lies in creating formulas ranging from foundational atomic propositions to intricate specifications.  
Leveraging the robust capabilities of the Z3 framework, the Formulizer seamlessly translates structural elements of the Kripke Structure into formal logical constructs. This meticulous process ensures the synthesis of expressive formulas that accurately capture system dynamics and properties.

*The Kripke Generator* is a pivotal software module tasked with the generation of Kripke structures. As its primary input, it requires a system model encapsulating robot entities. Initially, the module undertakes the process of system flattening, as described earlier, wherein the system's hierarchical representation is streamlined into a unified structure. Subsequently, leveraging the flattened model, the Kripke Generator proceeds to construct a Kripke structure. Additionally, this module assumes responsibility for the automated generation of systems based on specified input parameters, as detailed previously. The auto-generated system undergoes analogous flattening and subsequent conversion into a Kripke structure.

*The Solver* module assumes the critical role of executing the Bounded Model Checking (BMC) process. Its core function entails solving the formula generated by the Formulizer module iteratively within a specified bound until either a solution is found, no solution is possible, or a timeout threshold is surpassed. Furthermore, the Solver module encompasses the task of extracting the solution path from the formula's satisfying assignment, thereby providing insight into the trajectory followed to satisfy the specified constraints.

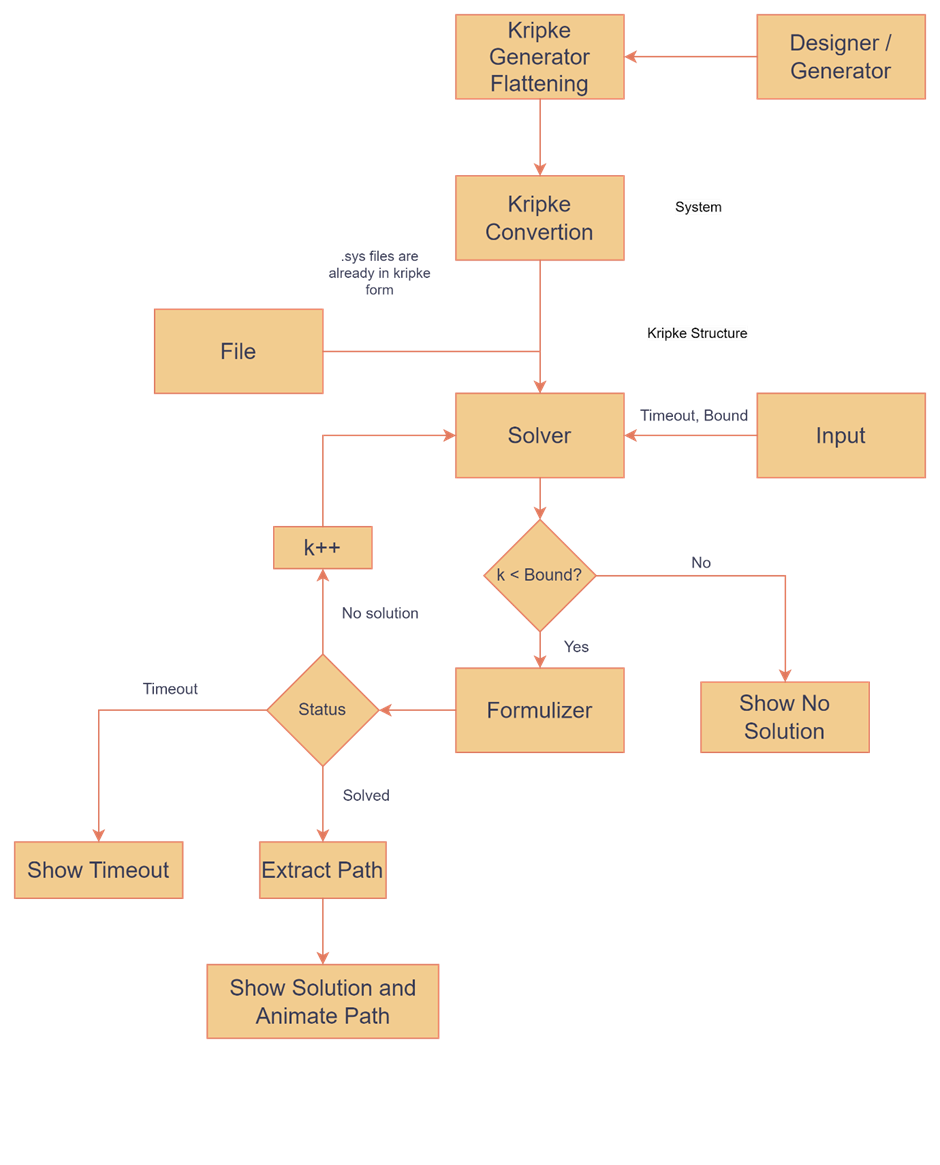


Figure 16 – System Process Flow

**Software Requirements and Specifications**

|  |  |  |
| --- | --- | --- |
| **Name** | **Description** | **Comply** |
| System Creation | The software can create systems in terms of grid and agents. | Yes |
| System Generation | The software can create systems automatically with certain constraints (grid size, counteragents count, etc..) | Yes |
| System Export/Import | The software can be saved as a file to the PC and can be loaded from a file. | Yes |
| Visualization | The software visualizes the loaded system as a Kripke structure.  The software visualizes the path finding process and algorithm advancement. | Yes |
| Algorithm | The software encodes the system as Boolean variables and use these variables to create the algorithm formulas | Yes |
| Solvers | The software attempts to satisfy the Boolean formula and the resulting assignment is decoded to the path taken. | Yes |

**Class Diagram**

Our project's class diagram is straightforward, serving to visually articulate the flow of data within our pipeline. As previously highlighted, this pipeline orchestrates a gradual metamorphosis of the data, transitioning it from a human-readable format to one that is computationally tractable.

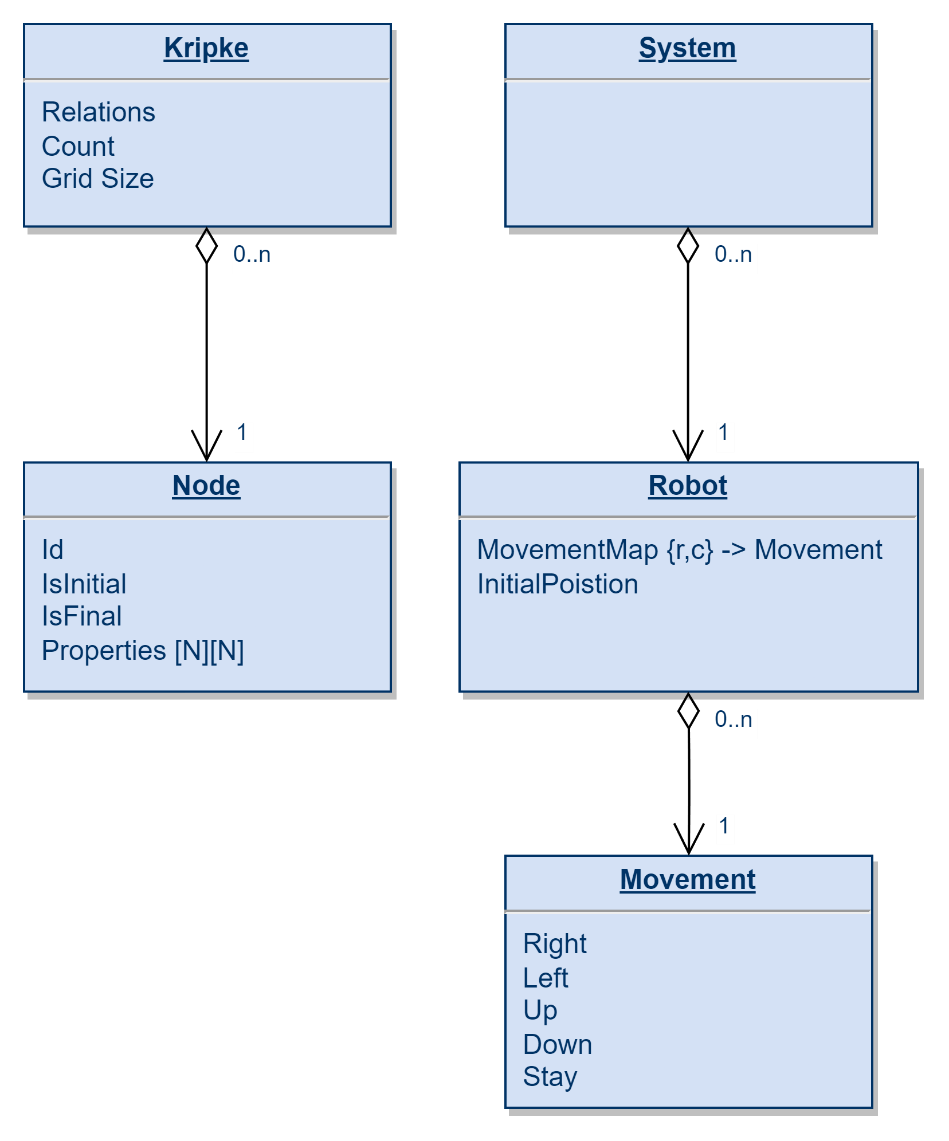


Figure 17 – Class Diagram

# 5. Analysis and Benchmarking

During our benchmarking and analysis phase, our primary objective was to identify the most influential factor in runtime. As previously mentioned, our parameters included grid size, number of counter agents, stray radius, stay chance and path length. To achieve this, we employed a methodical approach, systematically fixing parameters while adjusting a single variable at a time.

To ensure an accurate representation of computational performance, it was imperative to examine many samples for each variable.

Therefore, in these experiments every mark (blue dot) in the graph represents 100 individual tests.

For each experiment, a hypothesis was formulated and subsequently compared with the actual results obtained.

***The hypothesis is presented before each experiment shown below.***

# Number of counter agents

The hypothesis posited that augmenting the number of counter agents would lead to a corresponding increase in the count of potential occupants tracked during the flattening phase. However, it's important to note that the flattening phase serves as a preprocessing stage, consolidating multiple counter agents occupying the same position into a single occupancy indicator. Consequently, this process imposes a limitation on the total number of indicators within a single grid, constrained by the square of the grid cell count, which is .

Figure 18 – Counter Agents vs Solution Time,  
 Grid Size = 15, Radius = 10, Stay Chance = 0.1

The results indeed align with the hypothesis, demonstrating that varying the number of counter agents had only a negligible impact on solution time.

# Stray Radius

The hypothesis suggested that a larger stray radius would result in counter agents traversing the grid to a greater extent, surpassing their initial positions, and thereby increasing their contribution to grid occupancy. However, it was emphasized that this effect primarily influences the flattening phase, where all occupants within a cell are amalgamated into a single occupancy indicator.

Figure 19 – Stray Radius vs Solution Time,  
 Grid Size = 15, Counter = 3, Stay Chance = 0.1

The results are consistent with the hypothesis, revealing only a slight variation in solution time as the radius expands.

# Stay Chance

The hypothesis posited that increasing the stay chance of counter agents would heighten the probability of them staying in place, potentially resulting in prolonged occupancy of their positions. In addition, introducing an additional possibility of staying complicates subsequent movement predictions, as the counteragent may opt to remain stationary for an uncertain duration before transitioning to another cell. Nonetheless, it was observed that occupancy indicators remain bound by the dimensions of the grid.

Figure 20 – Stay Chance vs Solution Time,  
 Grid Size = 15, Counter = 3, Radius = 10

The results are mostly consistent with the hypothesis, revealing only a slight variation in solution time as the chance of having the option to remain stationary increases. However, a minor spike in runtime was always present when Stay chance was 0.4, the cause remains a mystery to us.

# Path Length

The hypothesis suggests that increasing the path length amplifies the number of steps an agent can undertake to discover an optimal route, thereby necessitating additional propositions in the formulas . The number of total atomic proposition (also known as base propositions) can be calculated by where is the path length and is the grid size. Consequently, this could extend the time required for the solver to find a satisfying assignment, potentially resulting in increased solution time.

Figure 21 – Path Length vs Solution Time,  
 Grid Size = 15, Counter = 3, Radius = 10, Stay Chance = 0.1

The results corroborate the hypothesis, demonstrating a consistent, nearly linear escalation as the path length increases. This indicates that the solver is compelled to exert more effort in finding an assignment that satisfies the formula under scrutiny.

# Grid Size ()

The hypothesis posits that enlarging the grid size results in a proportional increase in both the number of rows and columns, thereby initially generating more occupancy indicators. Crucially, this expansion also escalates the count of base propositions, which can be calculated as , where denotes the path length and denotes the grid size. As the quantity of base propositions swells, the length and intricacy of also expand, exerting a substantial influence on solution time and ultimately determining the algorithm's time complexity.

Figure 22 – Grid Size vs Solution Time,  
Counter = 3, Radius = 10, Stay Chance = 0.1

The findings are consistent with the hypothesis, showcasing a substantial surge in the volume of atomic propositions, expressions, and occupancy indicators. This surge triggers an exponential escalation in solution time, validating the anticipated correlation between the increased complexity of the problem and the corresponding rise in computational burden.

# Summary

The results shed light on the differentiation between three fundamental properties of grid-embedded systems: time, space, and density.

*Density* in this context refers to the concentration of counter agents on the grid and how they obstruct space within it. While the algorithm may not be directly influenced by density, the presence of blocked grids could indirectly lead to longer paths, subsequently elongating solution time.

*Time* in this context refers to the length of the path, as each step of the agent's movement increments this duration. With an increase in time, the solver is compelled to accommodate more constraints in its solution process, rendering the task of finding a satisfying assignment more challenging.

*Space* in this context refers to the spatial arrangement of the agent within the grid. As the grid size expands, the minimum length path () also increases. Consequently, the initial base proposition size becomes , resulting in a base size of

# Conclusions

The algorithm's objective is to reduce the grid pathfinding challenge into a SAT problem, achieved by translating it into boolean formulas representing space-time coordinates. This process harnesses solvers like Z3, which employ heuristic techniques and optimizations to tackle SAT problems effectively.

Given that the SAT problem is NP-complete, its response grows exponentially with input expansion. When transforming the pathfinding problem into SAT, the grid size emerges as the most influential input, followed by the path length.

# 6. Project Reflection

Initially, comprehending the topic proved to be a challenge, requiring extensive research and consultation with our supervisor. Despite encountering numerous obstacles, we persisted and eventually found effective solutions. Upon achieving a thorough understanding of the concept, we faced the challenge of translating it into machine-readable terms. It became apparent that the raw data (a grid) and the final processed data (the formula) posed difficulties in precise transformation. Subsequently, we devised a pipeline as our primary method of data processing.

This pipeline serves as a great example of the challenges we encountered. One of the initial hurdles was determining how to represent the grid and, more specifically, how to represent multiple counteragents within the same grid. Recognizing that the grid is an abstract concept with no direct representation in the code, we focused on the essential elements: the counteragents. To address this, we decomposed the problem into smaller components, creating objects of the type of Robot (counteragent). These objects encapsulated a data structure capable of mapping grid positions to potential movements. Instead of storing the entire grid, we only retained pertinent data, such as the grid cell occupied by the counteragent and its corresponding movements.

The second challenge proved to be more daunting as we encountered a major hurdle in transforming the grid structure into the Kripke structure. Initially, the system's complexity overwhelmed us as we attempted to meticulously track every movement of each counteragent, resulting in an overwhelming number of alternate realities. However, a breakthrough came when we realized that individual robot movements were not as crucial as identifying occupied grid cells. This insight allowed us to simplify the process by condensing multiple robots on the grid into a single representation with occupancy indicators. As a result, we were able to create a timeline illustrating the grid's occupancy state over time, effectively integrating the concept of time into our two-dimensional space resulting in solving a much more complicated problem than we intended.

The third challenge revolved around formulating the rules of the game. Despite having the timeline and understanding how the counteragents behaved, teaching our "good" robot the rules proved to be time-consuming. Initially, the robot's behavior was erratic: it would either teleport to the finish line or refuse to move altogether. At one point, it even duplicated itself across the grid. However, after extensive trial and error, we developed the formulas mentioned above, each designed to counteract the observed prohibited behaviors.

We credit ourselves for maintaining effective communication through regular meetings and brainstorming sessions, ensuring thorough planning before diving into development. Our approach emphasized sculpting our code and refining our algorithm to achieve optimal results. We utilized source control with GIT and SourceTree, adhering to industry standards by employing branching and pull requests as expected of software engineers. However, we acknowledge our shortcomings in GUI creation and design proficiency, though we made concerted efforts to produce the best possible outcomes in these areas.

# 9. References

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